

# SPLITTING BASED ITERATIVE METHODS FOR SOLVING LINEAR SYSTEMS

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# OUTLINE

- 1 SOLUTION OF LINEAR SYSTEMS
  - Iterative Methods
- 2 MATRIX SPLITTING BASED ITERATIVE SCHEMES
- 3 ALTERNATING ITERATIVE METHOD
  - Three step alternating iterative scheme
  - Preconditioned Iterative Method
  - Iterative scheme based on Regularization
- 4 NUMERICAL EXAMPLES
- 5 CONCLUSION AND REMARKS
- 6 REFERENCES

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# SYSTEM OF EQUATIONS

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**Note:** There are many problems which can be modelled as linear system equations.

# SYSTEM OF EQUATIONS

- A system of linear equations will be of the form

$$\begin{array}{cccccc} a_{11}x_1 & + & a_{12}x_2 & + & \cdots & + & a_{1n}x_n & = & b_1 \\ a_{21}x_1 & + & a_{22}x_2 & + & \cdots & + & a_{2n}x_n & = & b_2 \\ \vdots & & \vdots & & \ddots & & \vdots & & \vdots \\ a_{m1}x_1 & + & a_{m2}x_2 & + & \cdots & + & a_{mn}x_n & = & b_m \end{array}$$

- We can write in the matrix form as

$$\begin{pmatrix} a_{11} & a_{12} & \cdots & a_{1n} \\ a_{21} & a_{22} & \cdots & a_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ a_{m1} & a_{m2} & \cdots & a_{mn} \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{pmatrix} = \begin{pmatrix} b_1 \\ b_2 \\ \vdots \\ b_m \end{pmatrix}$$

- So the general form  $Ax = b$

# ELEMENTARY ROW OPERATIONS

The following operations applied to any matrix, yields a row-equivalent form.

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- **Scaling:** Multiplying a row by a nonzero constant ( $R_i \rightarrow cR_i$ ).
- **Replacement:** The row can be replaced by the sum of that row and a nonzero multiple of any other row ( $R_i \rightarrow R_i + cR_j$ ).

# SYSTEM OF LINEAR EQUATIONS

## An inconsistent example

Consider the following linear system

$$\begin{pmatrix} 1 & 2 \\ 2 & 4 \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \end{pmatrix} = \begin{pmatrix} 4 \\ 5 \end{pmatrix}$$

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Clearly this system of equations is not solvable or inconsistent.  
( Why?)

**No Solutions:** If  $\text{rank}(A) \neq \text{rank}([A|b])$ .

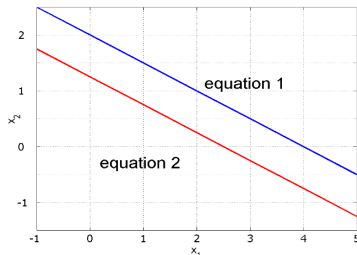


FIGURE: No solutions

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## Uniqueness of solutions

Consider the following linear system

$$\begin{pmatrix} 1 & 2 \\ 1 & -1 \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \end{pmatrix} = \begin{pmatrix} 4 \\ 2 \end{pmatrix}$$

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The above system has unique solution  $x_1 = 8/3$  and  $x_2 = 2/3$ .

**Unique Solution:** If  $\text{rank}(A) = \text{rank}([A|b]) = \text{No of variables}$ .

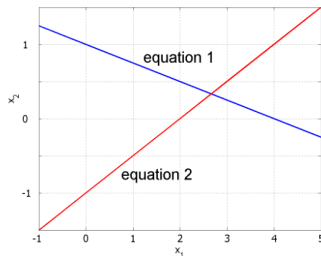


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## Rank deficient matrices/Infinite number of solutions

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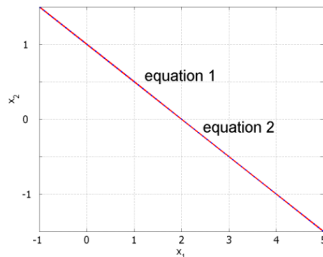


FIGURE: Unique Solution

**Unique Solution:** If  $\text{rank}(A) = \text{rank}([A|b]) < \text{No of variables}$ .



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### Direct Methods:

- Find a solution in a finite number of operations by transforming the system into an equivalent system that is 'easier' to solve.
- Diagonal, upper or lower triangular systems are easier to solve.
- Number of operations is a function of system size  $n$ .

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### Iterative Methods:

- Computes successive approximations of the solution vector  $x$  for a given  $A$  and  $b$ , starting from an initial point  $x_0$ .
- Total number of operations is uncertain, may not converge.

# GAUSSIAN ELIMINATION

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**Row Echelon Form:** After applying forward elimination, the augmented matrix will be in the following row echelon form:

$$\left[ \begin{array}{cccc|c} a'_{11} & a'_{12} & \cdots & a'_{1n} & b'_1 \\ 0 & a'_{22} & \cdots & a'_{2n} & b'_2 \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ 0 & 0 & \cdots & a'_{mn} & b'_m \end{array} \right]$$

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**Back Substitution:** Solve for  $x_i$ :

$$x_i = \frac{b'_i - \sum_{j=i+1}^n a'_{ij} x_j}{a'_{ii}}$$

# DRAWBACKS OF GAUSSIAN ELIMINATION

Recall:

**Example:** Solve the following system

$$1.133x_1 + 5.281x_2 = 6.414$$

$$24.14x_1 - 1.210x_2 = 22.93$$

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**Solution:** Using 4-digit rounding, we obtain the row echelon form as

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$$\left[ \begin{array}{cc|c} 1.133 & 5.281 & 6.414 \\ 0 & -113.7 & -113.8 \end{array} \right]$$

Hence  $x_1 = 0.9956$ ,  $x_2 = 1.001$  but the exact solution is  $x_1 = 1$ ,  $x_2 = 1$ .

# PIVOTING STRATEGIES (PARTIAL PIVOTING)

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- When the **pivotal element** is very small, the multipliers will be large.
- Adding numbers of widely different magnitudes can lead to a **loss of significance**.
- To reduce error, row interchanges are made to maximize the magnitude of the pivot element.

# PARTIAL PIVOTING

**Forward Elimination:** After applying forward elimination, the augmented matrix at step  $i$  will be in the following row echelon form:

$$\left[ \begin{array}{cccccccc|c} a_{11}^{(i)} & a_{12}^{(i)} & \cdots & a_{1i}^{(i)} & \cdots & a_{1j}^{(i)} & \cdots & a_{1n}^{(i)} & b_1^{(i)} \\ 0 & a_{22}^{(i)} & \cdots & a_{2i}^{(i)} & \cdots & a_{2j}^{(i)} & \cdots & a_{2n}^{(i)} & b_2^{(i)} \\ \vdots & \vdots & \ddots & \vdots & \ddots & \vdots & \ddots & \vdots & \vdots \\ 0 & 0 & \cdots & a_{ji}^{(i)} & \cdots & a_{jj}^{(i)} & \cdots & a_{jn}^{(i)} & b_i^{(i)} \\ \vdots & \vdots & \ddots & \vdots & \ddots & \vdots & \ddots & \vdots & \vdots \\ 0 & 0 & \cdots & a_{ji}^{(i)} & \cdots & a_{jj}^{(i)} & \cdots & a_{jn}^{(i)} & b_j^{(i)} \\ \vdots & \vdots & \ddots & \vdots & \ddots & \vdots & \ddots & \vdots & \vdots \\ 0 & 0 & \cdots & a_{mi}^{(i)} & \cdots & a_{mj}^{(i)} & \cdots & a_{mn}^{(i)} & b_m^{(i)} \end{array} \right]$$

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If  $\max\{|a_{ij}^{(i)}|, |a_{i+1i}^{(i)}|, \dots, |a_{ji}^{(i)}|, \dots, |a_{mi}^{(i)}|\} = |a_{ji}^{(i)}| \neq 0$  then swap row  $i$  with row  $j$ .

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Hence  $x_1 = 1$ ,  $x_2 = 1$ .

# DRAWBACKS OF PARTIAL PIVOTING

**Example:** Solve the following system

$$\begin{aligned}30x_1 + 591400x_2 &= 591700 \\5.291x_1 - 6.130x_2 &= 46.78\end{aligned}$$

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Hence  $x_1 = -10$ ,  $x_2 = 1.001$  but the exact solution is  $x_1 = 10$ ,  $x_2 = 1$ .

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- Condition number:  $k(A) = \|A\| \|A^{-1}\|$ .
- Effect of condition number:
- If the coefficient matrix is **ill-conditioned** then round-off will lead huge error in the solution.
- if  $A$  is ill-conditioned (a small change in some entries leads nonsingular to singular),  $A^{-1}$  will not be computed accurately.

# DRAWBACKS OF DIRECT METHODS

## Disasters due to bad numerics

On February 25, 1991, during the Gulf War, an American Patriot Missile battery in Dharan, Saudi Arabia, failed to track and intercept an incoming Iraqi Scud missile. This resulted in 28 deaths and 100 injuries.



# DRAWBACKS OF GAUSSIAN ELIMINATION

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On June 4, 1996 an unmanned Ariane 5 rocket launched by the European Space Agency exploded just 40 seconds after its lift-off from Kourou, French Guiana. Ariane explosion costing \$7 billion + The destroyed rocket and its cargo were valued at \$500 million.



and so on .....

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- For systems of small dimension, there is no need. Direct techniques will perform very efficiently.
- For systems with large, sparse coefficient matrices, direct techniques are often less efficient than iterative techniques.
- More appropriate when the number of equations involved is large, or when the matrix is sparse (many coefficients whose value is zero).



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## Iterative methods:

- Jacobi's Method (Carl Gustav Jakob Jacobi, 1804-1851)
- Gauss-Seidel Method (Carl Friedrich Gauss 1777-1855, Philipp Ludwig von Seidel 1821-1896)
- Successive Overrelaxation (SOR) Method

# JACOBI'S METHOD

Consider the linear system:

$$A\mathbf{x} = \mathbf{b}$$

where

$$A = \begin{pmatrix} a_{11} & a_{12} & \cdots & a_{1n} \\ a_{21} & a_{22} & \cdots & a_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ a_{n1} & a_{n2} & \cdots & a_{nn} \end{pmatrix}, \quad \mathbf{x} = \begin{pmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{pmatrix}, \quad \mathbf{b} = \begin{pmatrix} b_1 \\ b_2 \\ \vdots \\ b_n \end{pmatrix}.$$

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The Jacobi Method is derived by decomposing the matrix  $A$  into its diagonal ( $D$ ), strictly lower triangular ( $L$ ) and strictly upper triangular ( $U$ ) such that

$$A = D + L + U.$$

# JACOBI'S METHOD

- Thus the system  $Ax = b$  is rewritten as

$$x = -D^{-1}(L + U)x + D^{-1}b = Tx + c,$$

where  $T = -D^{-1}(L + U) = D^{-1}(D - A)$  and  $c = D^{-1}b$

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- Hence the Jacobi's iterative method in matrix form is given by

$$x^{(k)} = Tx^{(k-1)} + c, \text{ where } T = -D^{-1}(L + U), \text{ } c = D^{-1}b$$

and  $k = 1, 2, 3, \dots$

- Jacobi's iterative method in component form is given by

$$x_i^{(k)} = \frac{1}{a_{ii}} \left[ b_i - \sum_{\substack{j=1 \\ j \neq i}}^n a_{ij} x_j^{(k-1)} \right], \quad i = 1, 2, \dots, n \text{ and } k = 1, 2, 3, \dots$$

# JACOBI'S METHOD

**Example:** Use 5-digit rounding and Jacobi method to solve the following system:

$$\begin{aligned}2x_1 - x_2 &= 1 \\ -x_1 + 3x_2 - x_3 &= 8 \\ -x_2 + 2x_3 &= -5\end{aligned}$$



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**Solution:** From  $D = \begin{bmatrix} 2 & 0 & 0 \\ 0 & 3 & 0 \\ 0 & 0 & 2 \end{bmatrix}$ ,  $(L + U) = \begin{bmatrix} 0 & -1 & 0 \\ -1 & 0 & -1 \\ 0 & -1 & 0 \end{bmatrix}$ ,  $b = \begin{bmatrix} 1 \\ 8 \\ -5 \end{bmatrix}$ ,  
we have

$$T = -D^{-1}(L + U) = \begin{bmatrix} 0 & 0.5000 & 0 \\ 0.3333 & 0 & 0.3333 \\ 0 & 0.5000 & 0 \end{bmatrix}, \quad c = \begin{bmatrix} 0.5000 \\ 2.6667 \\ -2.5000 \end{bmatrix}$$

# JACOBI'S METHOD

**Iteration-1:** Taking  $x^{(0)} = [0 \ 0 \ 0]^T$ , we can compute

$$x^{(1)} = Tx^{(0)} + c = \begin{bmatrix} 0 \\ 0 \\ 0 \end{bmatrix} + \begin{bmatrix} 0.5000 \\ 2.6667 \\ -2.5000 \end{bmatrix} = \begin{bmatrix} 0.5000 \\ 2.6667 \\ -2.5000 \end{bmatrix}$$

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**Iteration-2:**

$$x^{(2)} = \begin{bmatrix} 0 & 0.5000 & 0 \\ 0.3333 & 0 & 0.3333 \\ 0 & 0.5000 & 0 \end{bmatrix} \begin{bmatrix} 0.5000 \\ 2.6667 \\ -2.5000 \end{bmatrix} + \begin{bmatrix} 0.5000 \\ 2.6667 \\ -2.5000 \end{bmatrix} = \begin{bmatrix} 1.8333 \\ 2.0000 \\ -1.1667 \end{bmatrix}$$

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$$x^{(1)} = Tx^{(0)} + c = \begin{bmatrix} 0 \\ 0 \\ 0 \end{bmatrix} + \begin{bmatrix} 0.5000 \\ 2.6667 \\ -2.5000 \end{bmatrix} = \begin{bmatrix} 0.5000 \\ 2.6667 \\ -2.5000 \end{bmatrix}$$

**Iteration-2:**

$$x^{(2)} = \begin{bmatrix} 0 & 0.5000 & 0 \\ 0.3333 & 0 & 0.3333 \\ 0 & 0.5000 & 0 \end{bmatrix} \begin{bmatrix} 0.5000 \\ 2.6667 \\ -2.5000 \end{bmatrix} + \begin{bmatrix} 0.5000 \\ 2.6667 \\ -2.5000 \end{bmatrix} = \begin{bmatrix} 1.8333 \\ 2.0000 \\ -1.1667 \end{bmatrix}$$

**After 20 iterations (Iteration-21):**

$$x^{(21)} = \begin{bmatrix} 2 \\ 3 \\ -1 \end{bmatrix} \Leftrightarrow x_1 = 2, x_2 = 3, x_3 = -1$$

# GAUSS-SEIDEL METHOD

- In Gauss-Seidel method, we rewrite the system  $Ax = b$  as

$$x = -(D + L)^{-1} Ux + (D + L)^{-1} b = Tx + c,$$

where  $T = -(D + L)^{-1} U$  and  $c = (D + L)^{-1} b$

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where  $T = -(D + L)^{-1} U$  and  $c = (D + L)^{-1} b$

- Hence the Gauss-Seidel iterative method in matrix form is given by

$$x^{(k)} = Tx^{(k-1)} + c, \text{ where } T = -(D + L)^{-1} U, \text{ } c = (D + L)^{-1} b$$

and  $k = 1, 2, 3, \dots$

- Jacobi's iterative method in component form is given by

$$x_i^{(k)} = \frac{1}{a_{ii}} \left[ b_i - \sum_{j=1}^{i-1} a_{ij} x_j^{(k)} - \sum_{j=i+1}^n a_{ij} x_j^{(k-1)} \right],$$

where  $i = 1, 2, \dots, n$  and  $k = 1, 2, 3, \dots$

# GAUSS-SEIDEL METHOD

**Example:** Use 5-digit rounding and Jacobi method to solve the following system:

$$\begin{aligned}2x_1 - x_2 &= 1 \\ -x_1 + 3x_2 - x_3 &= 8 \\ -x_2 + 2x_3 &= -5\end{aligned}$$

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**Solution:** From  $D + L = \begin{bmatrix} 2 & 0 & 0 \\ -1 & 3 & 0 \\ 0 & -1 & 2 \end{bmatrix}$ ,  $U = \begin{bmatrix} 0 & -1 & 0 \\ 0 & 0 & -1 \\ 0 & 0 & 0 \end{bmatrix}$ ,  $b = \begin{bmatrix} 1 \\ 8 \\ -5 \end{bmatrix}$ ,  
we have

$$T = -(D + L)^{-1}U = \begin{bmatrix} 0 & 0.5000 & 0 \\ 0 & 0.1667 & 0.3333 \\ 0 & 0.0833 & 0.1667 \end{bmatrix}, \quad c = \begin{bmatrix} 0.5000 \\ 2.8333 \\ -1.0833 \end{bmatrix}$$



# GAUSS-SEIDEL METHOD

**Iteration-1:** Taking  $x^{(0)} = [0 \ 0 \ 0]^T$ , we can compute

$$x^{(1)} = Tx^{(0)} + c = \begin{bmatrix} 0 \\ 0 \\ 0 \end{bmatrix} + \begin{bmatrix} 0.5000 \\ 2.8333 \\ -1.0833 \end{bmatrix} = \begin{bmatrix} 0.5000 \\ 2.8333 \\ -1.0833 \end{bmatrix}$$

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**Iteration-2:**

$$x^{(2)} = \begin{bmatrix} 0 & 0.5000 & 0 \\ 0 & 0.1667 & 0.3333 \\ 0 & 0.0833 & 0.1667 \end{bmatrix} \begin{bmatrix} 0.5000 \\ 2.8333 \\ -1.0833 \end{bmatrix} + \begin{bmatrix} 0.5000 \\ 2.8333 \\ -1.0833 \end{bmatrix} = \begin{bmatrix} 1.9167 \\ 2.9444 \\ -1.0278 \end{bmatrix}$$

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**Iteration-3:**

$$x^{(3)} = \begin{bmatrix} 0 & 0.5000 & 0 \\ 0 & 0.1667 & 0.3333 \\ 0 & 0.0833 & 0.1667 \end{bmatrix} \begin{bmatrix} 1.9167 \\ 2.9444 \\ -1.0278 \end{bmatrix} + \begin{bmatrix} 0.5000 \\ 2.8333 \\ -1.0833 \end{bmatrix} = \begin{bmatrix} 1.9722 \\ 2.9815 \\ -1.0093 \end{bmatrix}$$

# GAUSS-SEIDEL METHOD

Iteration-8:

$$x^{(8)} = \begin{bmatrix} 0 & 0.5000 & 0 \\ 0 & 0.1667 & 0.3333 \\ 0 & 0.0833 & 0.1667 \end{bmatrix} \begin{bmatrix} 1.9997 \\ 2.9998 \\ -1.0001 \end{bmatrix} + \begin{bmatrix} 0.5000 \\ 2.8333 \\ -1.0833 \end{bmatrix} = \begin{bmatrix} 1.9999 \\ 2.9999 \\ -1.0000 \end{bmatrix}$$

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Iteration-8:

$$x^{(8)} = \begin{bmatrix} 0 & 0.5000 & 0 \\ 0 & 0.1667 & 0.3333 \\ 0 & 0.0833 & 0.1667 \end{bmatrix} \begin{bmatrix} 1.9997 \\ 2.9998 \\ -1.0001 \end{bmatrix} + \begin{bmatrix} 0.5000 \\ 2.8333 \\ -1.0833 \end{bmatrix} = \begin{bmatrix} 1.9999 \\ 2.9999 \\ -1.0000 \end{bmatrix}$$

After 8 iterations (Iteration-9):

$$x^{(9)} = \begin{bmatrix} 0 & 0.5000 & 0 \\ 0 & 0.1667 & 0.3333 \\ 0 & 0.0833 & 0.1667 \end{bmatrix} \begin{bmatrix} 1.9999 \\ 2.9999 \\ -1.0000 \end{bmatrix} + \begin{bmatrix} 0.5000 \\ 2.8333 \\ -1.0833 \end{bmatrix} = \begin{bmatrix} 2 \\ 3 \\ -1 \end{bmatrix}$$

Therefore, we obtain the solution  $x_1 = 2$ ,  $x_2 = 3$ ,  $x_3 = -1$

# SPECTRAL RADIUS AND NORM OF A MATRIX

## Recall:

- Norm of a matrix ( $A_{n \times n}$ ):

$$\|A\|_{\infty} = \max_{1 \leq i \leq n} \sum_{j=1}^n |a_{ij}|, \|A\|_1 = \max_{1 \leq j \leq n} \sum_{i=1}^n |a_{ij}|, \|A\|_2 = \sqrt{\lambda_{\max}(A^*A)}, \|A\|_F = \sqrt{\sum_{i=1}^n \sum_{j=1}^n |a_{ij}|^2} = \sqrt{\text{trace}(A^*A)}.$$

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- **Spectral radius:** The spectral radius of a matrix  $A \in \mathbb{R}^{n \times n}$  is denoted by  $\rho(A)$  and defined by

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- Relation between norm and  $\rho(A)$ :  $\rho(A) \leq \|A\|$ .



# STOPPING CRITERIA

The following stopping rules are commonly used.

- Stop if the successive error,  $\|x^{(k)} - x^{(k-1)}\| < \epsilon$ .
- Stop if the residual error,  $\|b - Ax^{(k)}\| < \epsilon$ .
- Stop if the relative error,  $\frac{\|x^{(k)} - x^{(k-1)}\|}{\|x^{(k)}\|} < \epsilon$ .

# CONVERGENCE OF ITERATIVE METHODS

## THEOREM

*Consider a non-singular system  $Ax = b$  and its equivalent form be  $x = Tx + c$ . Then for any  $x^{(0)} \in \mathbb{R}^n$ , the sequence  $\{x^{(k)}\}$  defined by  $x^{(k)} = Tx^{(k-1)} + c$  ( $k = 1, 2, \dots$ ), converges to the unique solution  $A^{-1}b$ , of the system  $Ax = b$  if and only if  $\rho(T) < 1$ .*

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## COROLLARY

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# CONVERGENCE OF JACOBI'S AND GAUSS-SEIDEL METHOD

## THEOREM

If  $A$  is *strictly diagonally dominant* ( or *symmetric and positive definite* ), then for any choice  $x^{(0)}$ , the sequence  $\{x^{(k)}\}$  obtained by both the Jacobi and Gauss-Seidel iterative methods, converge to the unique solution  $A^{-1}b$ , of  $Ax = b$ .

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## THEOREM (STEIN-ROSENBERG)

If  $a_{ij} \leq 0$  for each  $i \neq j$  and  $a_{ii} > 0$  for each  $i = 1, 2, \dots, n$ . Then one and only one of the following statements holds:

- $0 \leq \rho(T_G) < \rho(T_J) < 1$
- $1 < \rho(T_J) < \rho(T_G)$
- $\rho(T_J) = 0 = \rho(T_G)$
- $\rho(T_J) = 1 = \rho(T_G)$

# SUCCESSIVE OVERRELAXATION (SOR) METHOD

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- Choose  $\omega$  such that  $\rho(T_{\omega}) < 1$  or  $\|T_{\omega}\| < 1$ .

# SOR METHOD BASED ON JACOBI METHOD

- It extends the Jacobi method by introducing a relaxation factor  $\omega$  to improve convergence.
- The iterative formula for the SOR method in matrix form, based on the Jacobi method, is:

$$x^{(k+1)} = (1 - \omega)x^{(k)} + \omega D^{-1} (b - (L + U)x^{(k)})$$

where  $A = D + L + U$  is the decomposition of matrix  $A$  into its diagonal ( $D$ ), strictly lower triangular ( $L$ ), and strictly upper triangular ( $U$ ).

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**Component-wise:** For  $i = 1, 2, \dots, n$ ,

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# SOR ITERATION METHODS

**Example:** Solve the following system

$$4x_1 + 3x_2 = 24, \quad 3x_1 + 4x_2 - x_3 = 30, \quad -x_2 + 4x_3 = -24$$

- by Gauss-Seidel method with  $x^{(0)} = (1, 1, 1)^T$
- by Gauss-Seidel with SOR and  $\omega = 1.25$ ,  $x^{(0)} = (1, 1, 1)^T$

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**Solution:**  $T_{gs} = \begin{bmatrix} 0 & -0.7500 & 0 \\ 0 & 0.5625 & 0.2500 \\ 0 & 0.1406 & 0.0625 \end{bmatrix}$ ,  $c_{gs} = \begin{bmatrix} 6.0000 \\ 3.0000 \\ -5.2500 \end{bmatrix}$

$$T_{\omega} = \begin{bmatrix} -0.2500 & -0.9375 & 0 \\ 0.2344 & 0.6289 & 0.3125 \\ 0.0732 & 0.1965 & -0.1523 \end{bmatrix}, \quad c_{\omega} = \begin{bmatrix} 7.5000 \\ 2.3438 \\ -6.7676 \end{bmatrix}$$

# SOR ITERATION METHODS

TABLE: Gauss-Seidel with SOR

$k$	1	2	3	4	5	6	7
$x^{(k)}$	6.3125	2.6223	3.1333	2.9571	3.0037	2.9963	3.0000
	3.5195	3.9585	4.0103	4.0075	4.0029	4.0009	4.0003
	-6.6501	-4.6004	-5.0967	-4.9735	-5.0057	-4.9983	-5.0003



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$k$	1	2	3	4	5	6	7
$x^{(k)}$	5.2500	3.1406	3.0879	3.0549	3.0343	3.0215	3.0134
	3.8125	3.8828	3.9268	3.9542	3.9714	3.9821	3.9888
	-5.0469	-5.0293	-5.0183	-5.0114	-5.0072	-5.0045	-5.0028

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## THEOREM (KAHAN)

Consider a system  $Ax = b$  where  $A \in \mathbb{R}^{n \times n}$  and  $b \in \mathbb{R}^n$ . If  $a_{ii} \neq 0$  for each  $i = 1, 2, \dots, n$  then  $\rho(T_\omega) \geq |\omega - 1|$ .

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**Note:** The above theorem tells us that the SOR method can converge only if  $0 < \omega < 2$ .

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**Note:** The above theorem tells us that the SOR method can converge only if  $0 < \omega < 2$ .

## THEOREM (OSTROWSKI-REICH)

Consider a system  $Ax = b$  where  $A \in \mathbb{R}^{n \times n}$  and  $b \in \mathbb{R}^n$ . If  $A$  is positive definite matrix and  $0 < \omega < 2$ , then the SOR method converges for any choice of  $x^{(0)}$ .

# OPTIMUM VALUE OF $\omega$ FOR SOR ITERATION METHODS

The optimal value of  $\omega$  minimizes the spectral radius of the iteration matrix.

## THEOREM

*Consider a system  $Ax = b$  where  $A \in \mathbb{R}^{n \times n}$  and  $b \in \mathbb{R}^n$ . If  $A$  is positive definite and tridiagonal, then*

- $\rho(T_{gs}) = \rho(T_j) < 1$
- *the optimal choice for  $\omega$  is*

$$\omega = \frac{2}{1 + \sqrt{1 - [\rho(T_j)]^2}},$$

*where  $T_j$  is the iteration matrix of Jacobi's Method.*

- $\rho(T_\omega) = \omega - 1$

# CLASSIFICATION OF SOR METHODS

**Under-relaxation:** If  $0 < \omega < 1$ , the method is called **under-relaxed**.

- It may converge slowly.
- It is used to make a non-convergent system converge, or to speedup convergence by avoiding oscillations





# OPTIMUM VALUE OF $\omega$ FOR SOR ITERATION METHODS

**Example:** Find the optimal choice of  $\omega$  for the following system

$$4x_1 + 3x_2 = 24, \quad 3x_1 + 4x_2 - x_3 = 30, \quad -x_2 + 4x_3 = -24$$

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**Solution:**  $T_j = -D^{-1}(L + U) = \begin{bmatrix} 0.25 & 0 & 0 \\ 0 & 0.25 & 0 \\ 0 & 0 & 0.25 \end{bmatrix} \begin{bmatrix} 0 & -3 & 0 \\ -3 & 0 & 1 \\ 0 & 1 & 0 \end{bmatrix}$

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$$= \begin{bmatrix} 0 & -0.75 & 0 \\ -0.75 & 0 & 0.25 \\ 0 & 0.25 & 0 \end{bmatrix}$$

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$$= \begin{bmatrix} 0 & -0.75 & 0 \\ -0.75 & 0 & 0.25 \\ 0 & 0.25 & 0 \end{bmatrix}$$

- Compute eigenvalues of  $T_j$  :  $\lambda_1 = 0$ ,  $\lambda_2 = \sqrt{0.625}$ ,  $\lambda_3 = -\sqrt{0.625}$ .

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**Example:** Find the optimal choice of  $\omega$  for the following system

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$$= \begin{bmatrix} 0 & -0.75 & 0 \\ -0.75 & 0 & 0.25 \\ 0 & 0.25 & 0 \end{bmatrix}$$

- Compute eigenvalues of  $T_j$  :  $\lambda_1 = 0$ ,  $\lambda_2 = \sqrt{0.625}$ ,  $\lambda_3 = -\sqrt{0.625}$ .
- Compute optimum value of  $\omega$  :

$$\omega = \frac{2}{1 + \sqrt{1 - 0.625}} \approx 1.24$$

## APPLICATION

Consider the two-dimensional Poisson's equation

$$-\frac{\partial^2 u}{\partial x^2} - \frac{\partial^2 u}{\partial y^2} = f(x, y), \quad (x, y) \in \blacksquare = [0, 1] \times [0, 1]$$

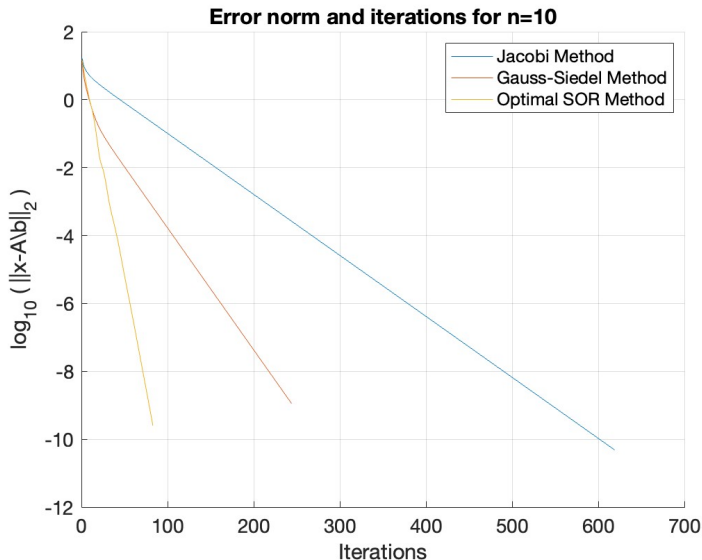
with  $u(x, y) = 0$  on the boundary  $\partial\blacksquare$ . Using 5-point stencil central difference scheme on a discretizing the unit square domain with  $n$  interior nodes, we obtain the following system

$$Ax = b, \quad A \in \mathbb{R}^{n^2 \times n^2}, \quad b \in \mathbb{R}^{n^2}$$

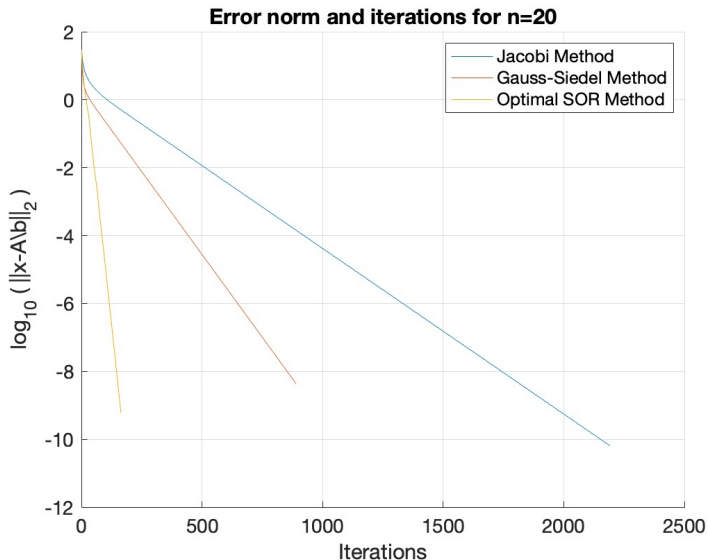
and the coefficient matrix will be of the form  $A = I \otimes P + P \otimes I$ , where

$$P = \begin{pmatrix} -2 & -1 & & 0 \\ -1 & \ddots & \ddots & \\ & \ddots & \ddots & -1 \\ 0 & & -1 & -2 \end{pmatrix}.$$

# APPLICATION

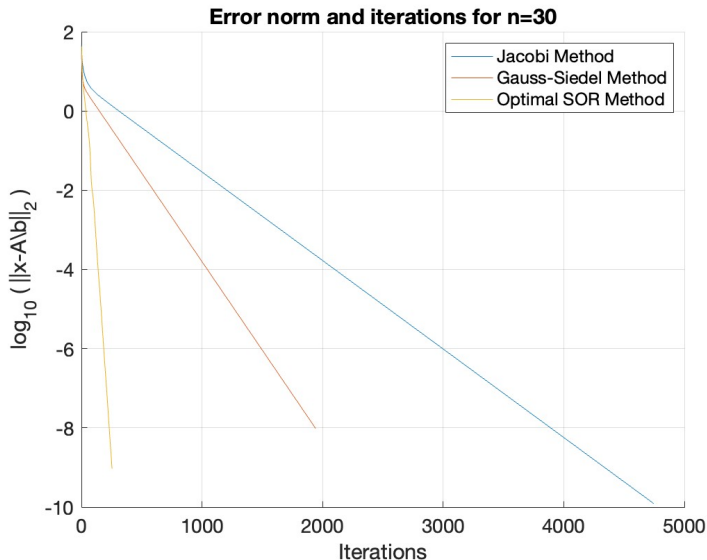


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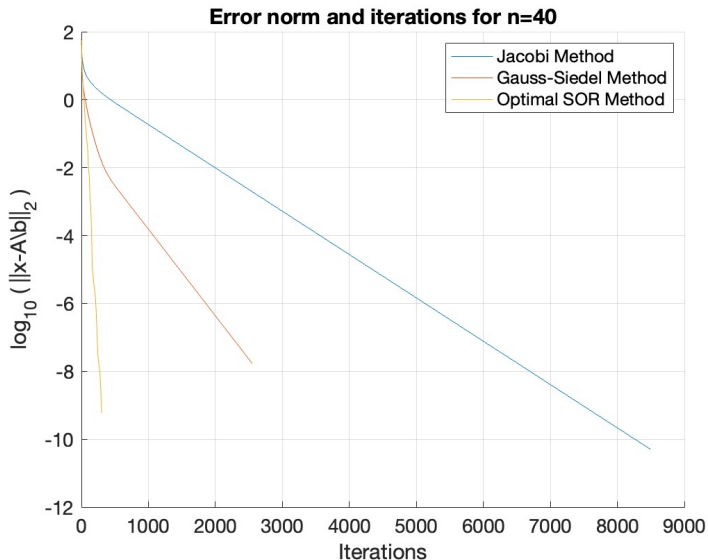




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# OUTLINE

- 1 SOLUTION OF LINEAR SYSTEMS
  - Iterative Methods
- 2 MATRIX SPLITTING BASED ITERATIVE SCHEMES
- 3 ALTERNATING ITERATIVE METHOD
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- A decomposition of the form  $A = B - C$  is called *splitting* of the matrix  $A \in \mathbb{R}^{m \times n}$
- To deal non singular system, several iterative methods are proposed to improve the convergence rate as well better complexity. For more details one can refer [3, 6, 7].



# MATRIX SPLITTING BASED ITERATIVE SCHEMES

- The singular and rectangular systems arise in various branches of Science and Engineering such as Markov Chain, Stochastic process, forecast modelling and partial differential equations.



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- The singular and rectangular systems arise in various branches of Science and Engineering such as Markov Chain, Stochastic process, forecast modelling and partial differential equations.
- To deal with such systems, in recent past many researchers have considered the splitting theory such as **proper regular splitting** and **proper weak regular** splitting. For example, If  $A = B - C$  is a proper splitting of  $A \in \mathbb{R}^{m \times n}$ , then the iterative scheme

$$x^{k+1} = B^\dagger Cx^k + B^\dagger b \quad (1)$$

for (1) converges to  $A^\dagger b$ , **the least squares solution** for any initial vector  $x^0$  iff  $\rho(B^\dagger C) < 1$ .

# INTRODUCTION

## Drawbacks:

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In this talk, we will discuss an alternating iterative scheme which can avoid the spectral radius calculation.

# BASIC TERMINOLOGY

## DEFINITION (MOORE-PENROSE INVERSE)

Let  $A \in \mathbb{C}^{m \times n}$ . If a matrix  $X \in \mathbb{C}^{n \times m}$  satisfies the following properties  $AXA = A$ ,  $XAX = X$ ,  $(AX)^* = AX$ ,  $(XA)^* = XA$ , then  $X$  is called the **Moore-Penrose inverse** of  $A$  and denoted as  $A^\dagger$ .

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[5] L. Jena, D. Mishra, and S. Pani. Convergence and comparison theorems for single and double decompositions of rectangular matrices. *Calcolo*, 51(1):141-149, 2014.



# BASIC TERMINOLOGY & RESULTS

Based on the above definitions and splitting, the following results have been proved in [5] and [2].

## THEOREM (THEOREM 1.3, [5])

*Let  $A = B - C$  be a proper regular splitting of  $A \in \mathbb{R}^{m \times n}$ . Then the Moore-Penrose inverse  $A^\dagger \geq 0$  if and only if the spectral radius of the iteration matrix is less than 1, i.e.,  $\rho(B^\dagger C) < 1$ .*

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[2] A. Berman and R. J. Plemmons. Cones and iterative methods for best least squares solutions of linear systems. SIAM Journal on Numerical Analysis, 11(1):145-154, 1974.

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## THREE STEP ALTERNATING ITERATIVE SCHEME

- Let  $A = B - C = X - Y = S - T$  be three proper splittings of the matrix  $A \in \mathbb{R}^{m \times n}$ . The followings are the proposed iterative schemes for the above three splittings

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which provide the solution of the system (1) iteratively for any initial guess  $x^0$ .

- By simplifying the iterative schemes (2), (3) and (4) we have the alternating iteration

$$x^{k+1} = Hx^k + Qb, \quad (5)$$

where  $H = S^\dagger TX^\dagger YB^\dagger C$  and  $Q = S^\dagger(TX^\dagger YB^\dagger + TX^\dagger + I)$ .



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Note that convergence of individual splitting does not imply the convergence of alternating iterative scheme which can be seen in the next example.

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### EXAMPLE

$$\text{Consider } A = \begin{bmatrix} 1 & 5 & -2 & -3 \\ 2 & -2 & 4 & -2 \\ -1 & 0 & -1 & 3 \end{bmatrix} = \begin{bmatrix} \frac{33}{10} & \frac{-21}{10} & \frac{9}{2} & \frac{-51}{5} \\ \frac{-11}{5} & \frac{32}{5} & \frac{-11}{2} & \frac{93}{10} \\ \frac{-23}{5} & \frac{77}{10} & \frac{-48}{5} & \frac{119}{10} \end{bmatrix} - \begin{bmatrix} \frac{23}{10} & \frac{-71}{10} & \frac{13}{2} & \frac{-36}{5} \\ \frac{-21}{5} & \frac{42}{5} & \frac{-19}{2} & \frac{113}{10} \\ \frac{-18}{5} & \frac{77}{10} & \frac{-43}{5} & \frac{89}{10} \end{bmatrix}$$

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 &= D - E = F - G = J - K
 \end{aligned}$$

are three proper splittings of  $A$ . Here  $\rho(D^\dagger E) = 0.5899 < 1$ ,  $\rho(F^\dagger G) = 0.8378 < 1$ ,  $\rho(J^\dagger K) = 0.8713 < 1$  but  $\rho(H) = 2.1125 \not< 1$ .

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## THEOREM (1)

*Let  $A = B - C = X - Y = S - T$  be three proper regular splittings of a semi-monotone matrix  $A \in \mathbb{R}^{m \times n}$ . Then  $\rho(H) = \rho(S^\dagger TX^\dagger YB^\dagger C) < 1$ .*

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## THEOREM (2)

*Let  $A \in \mathbb{R}^{m \times n}$  be a semi-monotone matrix and  $A = D - E = F - G = J - K$  be three proper regular splittings of  $A$  with  $R(D + J - A + KF^\dagger E) = R(A)$  and  $N(D + J - A + KF^\dagger E) = N(A)$ . Then,  $\rho(H) \leq \min\{\rho(D^\dagger E), \rho(F^\dagger G), \rho(J^\dagger K)\} < 1$ .*

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where the matrix  $P$  is a non singular matrix of order  $m$ . Let  $PA = K_p - L_p$  be a splitting of the matrix  $PA \in \mathbb{R}^{n \times n}$ , where  $K_p$  and  $L_p$  have same order as of  $PA$ .

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If  $PA = K_p - L_p$  is a proper splitting of the matrix  $PA$ , then the iterative scheme (7) will converge to the least square solution  $A^\dagger b$  for any initial guess  $x^0$  if and only if  $\rho(K_p^\dagger L_p) < 1$ .

# PRECONDITIONED ITERATIVE METHOD

The comparison between preconditioned approach and proper weak regular splitting approach has discussed in the next theorem.

## THEOREM

*Let  $A = M - N$  be a proper regular splitting of a semi-monotone matrix  $A \in \mathbb{R}^{m \times n}$ . Assume that there exists an orthogonal matrix  $P \in \mathbb{R}^{m \times m}$  such that  $A^\dagger P^{-1} \geq 0$ . If  $PA = M_p - N_p$  is a proper weak regular splitting of  $PA$  and  $M_p^\dagger P \geq M^\dagger$ , then  $\rho(M_p^\dagger N_p) \leq \rho(M^\dagger N) < 1$ .*

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### THEOREM

*Let  $A = M - N$  be a convergent proper splitting of  $A \in \mathbb{R}^{m \times n}$ . Let  $P \in \mathbb{R}^{m \times m}$  be a nonpositive orthogonal matrix such that  $PA = M_p - N_p$  is a proper regular splitting of  $PA$ . If  $A^\dagger \leq 0$  and  $M_p^\dagger P \leq M^\dagger$ , then  $\rho(M_p^\dagger N_p) \leq \rho(M^\dagger N) < 1$ .*





# ITERATIVE SCHEME BASED ON REGULARIZATION

- To obtain the unique least square solution  $A^\dagger b$  of  $Ax = b$ , we need to solve the normal equation  $A^T Ax = A^T b$ . But in general the matrix  $A^T A$  is ill-conditioned matrix [4].
- Therefore, we consider the following well-posed linear system:

$$(A^T A + \lambda I)x = A^T b, \quad (8)$$

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where  $\lambda > 0$  is called regularization parameter.

- In [1], it is proved that the matrix  $A^T A + \lambda I$  is nonsingular for every  $\lambda > 0$ . If we assume  $B = A^T A + \lambda I$ , then the system (8) reduces to the following nonsingular system:

$$Bx = A^T b. \quad (9)$$

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- We consider  $B = M_\lambda - N_\lambda$  is a splitting of the nonsingular matrix  $B$ , then the iterative scheme

$$x^{k+1} = M_\lambda^{-1}N_\lambda x^k + M_\lambda^{-1}A^Tb \quad (10)$$

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We have the following comparison result for the above setting:

## THEOREM

*Suppose  $A = M - N$  be a proper convergent weak splitting of type II. Let  $B = M_\lambda - N_\lambda$  be a convergent weak splitting of type I of the matrix  $B$ . If  $M_\lambda^{-1} A^T \geq M^\dagger$ , then  $\rho(M_\lambda^{-1} N_\lambda) \leq \rho(M^\dagger N) < 1$ .*

# OUTLINE

- 1 SOLUTION OF LINEAR SYSTEMS
  - Iterative Methods
- 2 MATRIX SPLITTING BASED ITERATIVE SCHEMES
- 3 ALTERNATING ITERATIVE METHOD
  - Three step alternating iterative scheme
  - Preconditioned Iterative Method
  - Iterative scheme based on Regularization
- 4 NUMERICAL EXAMPLES
- 5 CONCLUSION AND REMARKS
- 6 REFERENCES

# NUMERICAL EXAMPLES

## EXAMPLE (1)

Consider the system  $Ax = b$ , where  $A = \begin{bmatrix} 6.2 & 9.7 & -7.5 & -4.3 \\ 3.4 & -8.8 & 2.6 & 5.0 \\ -7.3 & -2.8 & 6.1 & 1.3 \end{bmatrix}$

and  $b = (0, 1, -1)^T$ . Clearly  $A$  is semi-monotone matrix since  $A^\dagger \geq 0$ .

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and  $b = (0, 1, -1)^T$ . Clearly  $A$  is semi-monotone matrix since  $A^\dagger \geq 0$ . Consider the following three proper regular splittings of  $A$  as

$$A = \begin{bmatrix} 7.39962 & 9.96576 & -7.26446 & -3.80128 \\ 5.35168 & -7.64282 & 3.33142 & 5.80321 \\ -6.19108 & -1.23196 & 7.16198 & 1.85758 \end{bmatrix} - \begin{bmatrix} 1.19962 & 0.26576 & 0.23554 & 0.498716 \\ 1.95168 & 1.15718 & 0.73142 & 0.803207 \\ 1.10892 & 1.56804 & 1.06198 & 0.557576 \end{bmatrix} \quad (\text{Spl-1})$$

$$= \begin{bmatrix} 6.76084 & 11.3339 & -6.45554 & -3.97232 \\ 5.8871 & -8.76498 & 3.64066 & 6.46661 \\ -6.45554 & -1.09718 & 7.48736 & 1.89716 \end{bmatrix} - \begin{bmatrix} 0.56084 & 1.63388 & 1.04446 & 0.327683 \\ 2.4871 & 0.03502 & 1.04066 & 1.46661 \\ 0.84446 & 1.70282 & 1.38736 & 0.597162 \end{bmatrix} \quad (\text{Spl-2})$$

$$= \begin{bmatrix} 7.48756 & 10.7622 & -6.48554 & -3.56629 \\ 6.09602 & -8.5369 & 3.79814 & 6.56695 \\ -6.28704 & -2.2252 & 6.56918 & 1.77155 \end{bmatrix} - \begin{bmatrix} 1.28756 & 1.06216 & 1.01446 & 0.733707 \\ 2.69602 & 0.2631 & 1.19814 & 1.56695 \\ 1.01296 & 0.5748 & 0.46918 & 0.471554 \end{bmatrix} \quad (\text{Spl-3})$$



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Here the spectral radius of the iteration matrix, i.e.,  $\rho(H) = 0.492122$  which is less than  $\min\{\rho(K^\dagger L) = 0.774462, \rho(U^\dagger V) = 0.8130275, \rho(X^\dagger Y) = 0.787961\} < 1$ .

# NUMERICAL EXAMPLES

## EXAMPLE (2)

$$\text{Let } A = \begin{bmatrix} 8.3 & -6.7 & 4.0 & -2.6 \\ -7.0 & 2.9 & 0.9 & -1.3 \\ 7.7 & -3.2 & -7.4 & 3.1 \end{bmatrix} = K - L$$

$$= \begin{bmatrix} -1.49724 & -13.2765 & 11.0203 & -15.1059 \\ -7.3728 & -1.87972 & -3.13564 & -4.5578 \\ 13.4498 & -3.50256 & -12.5629 & 7.12079 \end{bmatrix} - \begin{bmatrix} -9.79724 & -6.57652 & 7.02032 & -12.5059 \\ -0.3728 & -4.77972 & -4.03564 & -3.2578 \\ 5.7498 & -0.30256 & -5.16292 & 4.02079 \end{bmatrix}$$

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$$\text{Let } P = \begin{bmatrix} -1 & 0 & 0 \\ 0 & 0 & -1 \\ 0 & -1 & 0 \end{bmatrix} \text{ be an nonpositive orthogonal matrix such that } PA = K_P - L_P$$

$$= \begin{bmatrix} -7.09049 & 8.31709 & -2.38373 & 4.18852 \\ -7.39842 & 3.63104 & 8.74419 & -2.92637 \\ 7.98315 & -2.80637 & 0.346307 & 1.55539 \end{bmatrix} - \begin{bmatrix} -7.09049 & 8.31709 & -2.38373 & 4.18852 \\ -7.39842 & 3.63104 & 8.74419 & -2.92637 \\ 7.98315 & -2.80637 & 0.346307 & 1.55539 \end{bmatrix}.$$

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is a proper regular splitting of  $PA$ .

Here  $A^\dagger \leq 0$ ,  $\rho(K_p^\dagger L_p) = 0.5989862 \leq 0.8617974 = \rho(K^\dagger L) < 1$ .

# NUMERICAL EXAMPLES

TABLE: Convergence Analysis of Alternating Scheme

Example	$\epsilon$	N	$\ Ax_k - b\ _2$	$\ x_k - A^\dagger b\ _2$	MT
Ex-1	$10^{-9}$	8	$2.8929e^{-11}$	$6.2249e^{-12}$	0.0015
Ex-1	$10^{-15}$	12	$1.4315e^{-16}$	$3.0803e^{-12}$	0.0024

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TABLE: Comparison Analysis of Alternating Scheme

Splittings	$\epsilon$	N	$\ Ax_k - b\ _2$	$\ x_k - A^\dagger b\ _2$	MT
Alternating Scheme	$10^{-9}$	8	$2.8929e^{-11}$	$6.2249e^{-12}$	0.0015
Splitting 1	$10^{-9}$	14	$1.5582e^{-9}$	$2.9480e^{-10}$	0.0026
Splitting 2	$10^{-9}$	22	$1.9168e^{-9}$	$4.2316e^{-10}$	0.0039
Splitting 3	$10^{-9}$	28	$2.9515e^{-9}$	$6.2998e^{-10}$	0.0049

# NUMERICAL EXAMPLES

TABLE: Comparison Analysis between  $K^\dagger L$  and  $K_p^\dagger L_p$

Splittings	$\epsilon$	N	$\ Ax_k - b\ _2$	$\ x_k - A^\dagger b\ _2$	MT
$K - L$	$10^{-9}$	204	$2.2831e^{-7}$	$9.3700e^{-9}$	0.1361
$K_p - L_p$	$10^{-9}$	32	$1.1962e^{-8}$	$6.1307e^{-10}$	0.0073
$K - L$	$10^{-15}$	345	$2.1527e^{-13}$	$8.8351e^{-15}$	0.3721
$K_p - L_p$	$10^{-15}$	51	$1.4127e^{-14}$	$7.24056e^{-16}$	0.0133

# NUMERICAL EXAMPLES

## EXAMPLE (3)

Let us consider the following two-dimensional partial differential equation

$$-\frac{\partial^2 u}{\partial^2 x} - \frac{\partial^2 u}{\partial^2 y} + 0.5 \frac{\partial u}{\partial x} + 2 \frac{\partial u}{\partial y} = f(x, y), \quad (x, y) \in [0, 1] \times [0, 1]$$

with Dirichlet boundary conditions.



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with Dirichlet boundary conditions. If we use central difference scheme on a uniform grid with  $(N + 2)$  nodes, then we will obtain a linear system  $Ax = b$ , where the coefficient matrix is of order  $N^2$  and of the following form

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with Dirichlet boundary conditions. If we use central difference scheme on a uniform grid with  $(N + 2)$  nodes, then we will obtain a linear system  $Ax = b$ , where the coefficient matrix is of order  $N^2$  and of the following form

$$A = I \otimes P + Q \otimes I,$$

$$P = \text{trid}(-(h+1), 4, (h-1)) \text{ and } Q = \text{trid}\left(-\frac{h+4}{4}, 0, \frac{h-4}{4}\right).$$

# NUMERICAL EXAMPLES

The comparison analysis of the three step with the schemes of [8], and [9] are summarized in Table 6.

- [8] S.Q. Shen and T. Z. Huang. Convergence and comparison theorems for double splittings of matrices. Comput. Math. Appl. 51(12):1751-1760, 2006.
- [9] S. Srivastava, D. Gupta, and A. Singh. An iterative method for solving singular linear systems with index one. Afrika Matematika, 27(5-6):815-824, 2016

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TABLE: Comparison of error bounds and mean processing time for  $\epsilon = 10^{-12}$

Order of $A$	Method	$\ Ax_k - b\ _2$	$\ x_k - A^{-1}b\ _2$	MT
100 ( $N = 10$ )	Method of [9]	$8.6544e^{-15}$	$1.5445e^{-14}$	0.00845
	Method of [8]	$3.7380e^{-12}$	$2.1915e^{-11}$	0.00138
	Three-step	$1.8618e^{-13}$	$5.2171e^{-13}$	0.00055
400 ( $N = 20$ )	Method of [9]	$4.3592e^{-14}$	$8.0318e^{-14}$	0.27489
	Method of [8]	$3.8949e^{-12}$	$8.2797e^{-11}$	0.37036
	Three-step	$2.0225e^{-13}$	$1.5958e^{-12}$	0.00329
900 ( $N = 30$ )	Method of [9]	$1.4683e^{-13}$	$3.9423e^{-13}$	3.23154
	Method of [8]	$3.9572e^{-12}$	$1.8246e^{-10}$	4.65695
	Three-step	$5.6277e^{-13}$	$3.0136e^{-12}$	0.0472
1600 ( $N = 40$ )	Method of [9]	$3.2206e^{-13}$	$3.0576e^{-12}$	18.58005
	Method of [8]	$4.0005e^{-12}$	$3.2221e^{-10}$	23.28531
	Three-step	$1.3450e^{-12}$	$4.1725e^{-12}$	0.19843

# NUMERICAL EXAMPLES

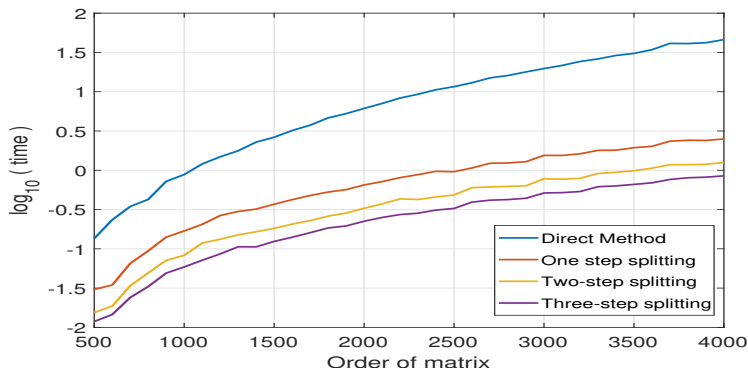


FIGURE: Comparison of mean processing time for different order matrices

# NUMERICAL EXAMPLES

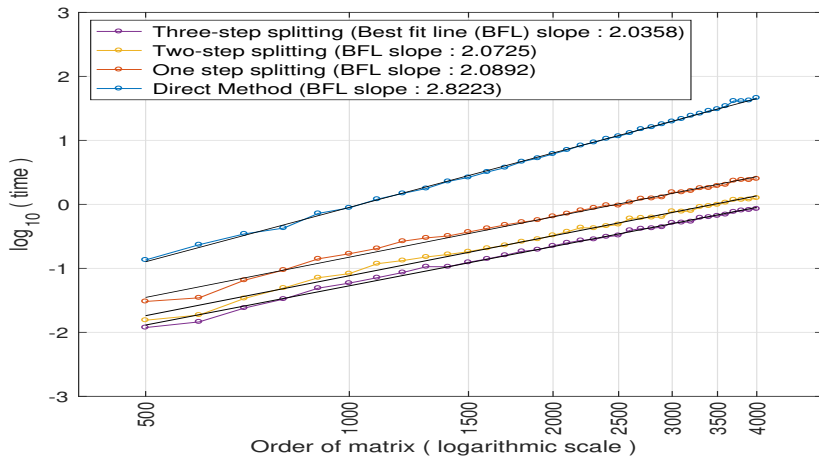


FIGURE: Comparison of time complexity for different order matrices

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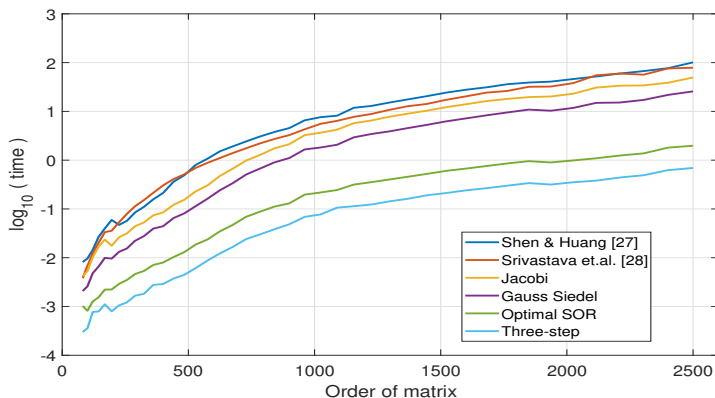


FIGURE: Comparison of mean processing time with existing methods

# NUMERICAL EXAMPLES

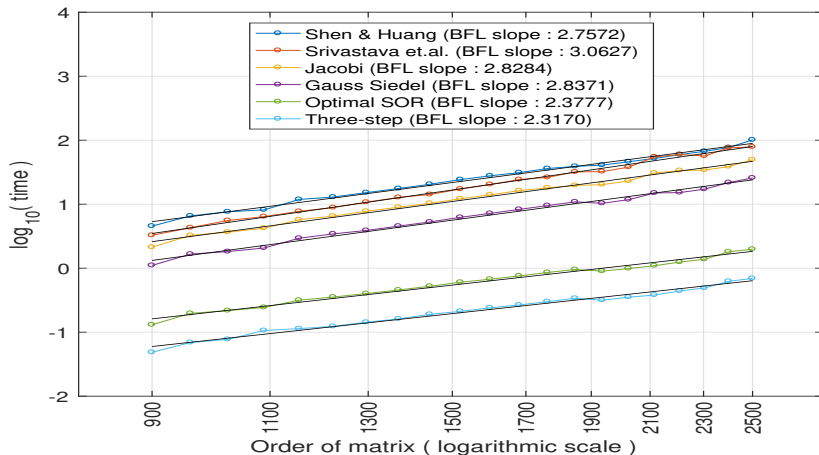


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# CONCLUSION AND REMARKS

- We have discussed theoretical results for the proposed alternating iterative schemes. Numerical examples are provided to justify the schemes.
- The proposed scheme converges much faster than the well-known splittings. We also discuss a suitable choice of preconditioned matrix, or regularization parameter can make the system well-posed.
- Other regularization techniques can be further used for solving singular system.

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Any Questions/comments ?